

Finding the MEK_L Performance

Definitions

Definitions

```
ID = {{1, 0}, {0, 1}};  
X = {{0, 1}, {1, 0}};  
Z = {{1, 0}, {0, -1}};  
Y = i * X.Z;  
HL = Cos[θ] * X + Sin[θ] Z;  
HLη[xη_] = HL.MatrixPower[i * Y, xη];
```

Defining some matrices

```
ρONE[ε_] = (ID + (1 - 2 ε) * HL) / 2;  
ρPerfectGlobal = KroneckerProduct[(ID + HL) / 2, (ID + HL) / 2];  
ρPerfectLocal1 = KroneckerProduct[ID, (ID + HL) / 2];  
ρPerfectLocal2 = KroneckerProduct[ID, (ID + HL) / 2];  
ρ[ε_] = KroneckerProduct[ρONE[ε], ρONE[ε];
```

We think of MEK being a protocol that takes 2 magic states in the state $\rho[\epsilon]$ and pumps on them with a circuit using a further 8 copies.

Channels

Channels

```
U1[x1_, x2_, x3_, x4_, x5_, x6_, x7_, x8_] =  
  MatrixPower[KroneckerProduct[X, X], x1 + x3].  
  MatrixPower[KroneckerProduct[Z, Z], x2 + x4];  
U2[x1_, x2_, x3_, x4_, x5_, x6_, x7_, x8_] =  
  MatrixPower[KroneckerProduct[X, X], x5 + x7].  
  MatrixPower[KroneckerProduct[Z, Z], x6 + x8];  
Proj[x1_, x2_, x3_, x4_, x5_, x6_, x7_, x8_, xη_] =  
  KroneckerProduct[ID, HLη[xη]] . (KroneckerProduct[ID, ID] +  
    (-1)^(x3+x4+x5+x6) KroneckerProduct[HLη[xη], HLη[xη]]) / 2;  
  
Kraus[x1_, x2_, x3_, x4_, x5_, x6_, x7_, x8_, xη_] :=  
  U2[x1, x2, x3, x4, x5, x6, x7, x8].  
  Proj[x1, x2, x3, x4, x5, x6, x7, x8, xη].U1[x1, x2, x3, x4, x5, x6, x7, x8];
```

We will apply a noisy channel to $\rho[\epsilon]$ which is a sum over Kraus operators labelled by $[x1_, x2_, x3_, x4_, x5_, x6_, x7_, x8_]$, where $x_j=1$ indicates an error on the j^{th} of the 8 pumping magic states.

For each Kraus operator we further break it up into 3 parts U1, Proj and U2.

U1 and U2 are unitaries that consist of Pauli errors.

Proj is a noisy parity projection, which is usually a projection onto the even parity Hadamard space, though some Pauli errors can generate a projection onto the odd parity Hadamard space.

```
Channel[x1_, x2_, x3_, x4_, x5_, x6_, x7_, x8_, xη_, ρin_] :=
  Kraus[x1, x2, x3, x4, x5, x6, x7, x8, xη].ρin.
  ConjugateTranspose[Kraus[x1, x2, x3, x4, x5, x6, x7, x8, xη]]
```

The above code simply takes a Kraus operator and makes a quantum channel from it

```
w0 = Permutations[{0, 0, 0, 0, 0, 0, 0, 0}];
w2 = Permutations[{1, 1, 0, 0, 0, 0, 0, 0}];
w4 = Permutations[{1, 1, 1, 1, 0, 0, 0, 0}];
w6 = Permutations[{1, 1, 1, 1, 1, 1, 0, 0}];
w8 = {{1, 1, 1, 1, 1, 1, 1, 1}};
w = Union[w0, w2, w4, w6, w8];
size = Dimensions[w][[1]];
```

We construct a list of all nontrivial even weight errors (odd weight errors are always detected by the circuit)

```
FullChannel[ε_, δ_, ρin_] :=
  ((1 - δ) * Sum[(ε / (1 - ε))^Total[w[[k]]] * (1 - ε)^8 * Channel[w[[k, 1]], w[[k, 2]], w[[k, 3]],
    w[[k, 4]], w[[k, 5]], w[[k, 6]], w[[k, 7]], w[[k, 8]], 0, ρin],
    {k, 1, size}]) + (δ * Sum[(ε / (1 - ε))^Total[w[[k]]] * (1 - ε)^8 *
    Channel[w[[k, 1]], w[[k, 2]], w[[k, 3]], w[[k, 4]], w[[k, 5]],
    w[[k, 6]], w[[k, 7]], w[[k, 8]], 1, ρin], {k, 1, size}]);
```

Above we form a channel by summing over all errors of even weight (those not detected) and weight them by $\epsilon^N * (1 - \epsilon)^{N-8}$ where N is the number of errors (the number of nonzero values for [x1_,x2_,x3_,x4_,x5_,x6_,x7_,x8_],)

Output state

```
ρOUT[ε3_, εk_, δ_] = FullChannel[ε3, δ, ρ[εk]];
```

We find $\rho_{OUT}[\epsilon]$ by applying the noisy map $\text{FullChannel}[\epsilon, \rho_{in}]$ to the input $\rho[\epsilon]$. Remember the input consists of 2 noisy magic states and 8 noisy magic were responsible for implementing the noisy map.

Polynomials

Polynomials

```

Psuc[ε3_, ek_, δ_] =
  Simplify[Simplify[Simplify[Tr[ρOUT[ε3, ek, δ] /. {Conjugate[Sin[θ]] → Sin[θ],
    Conjugate[Cos[θ]] → Cos[θ]}] /.
    {Conjugate[Sin[θ]] → Sin[θ], Conjugate[Cos[θ]] → Cos[θ]}]] /.
    {Conjugate[Sin[θ]] → Sin[θ], Conjugate[Cos[θ]] → Cos[θ]}]
1 - 448 ε35 + 448 ε36 - 256 ε37 + 64 ε38 -  $\frac{1}{2} \delta (1 - 2 \epsilon 3)^4 (1 - 2 \epsilon k)^2 - 2 \epsilon k + 2 \epsilon k^2 -$ 
  64 ε33 (2 - εk + εk2) + 32 ε34 (9 - εk + εk2) - 8 ε3 (1 - 2 εk + 2 εk2) + 8 ε32 (5 - 6 εk + 6 εk2)

```

The success probability.

```

εout[ε3_, ek_, δ_] =  $\frac{1}{\text{Psuc}[\epsilon 3, \epsilon k, \delta]}$  * Expand[Simplify[
  Simplify[Tr[ρOUT[ε3, ek, δ].(KroneckerProduct[ID, ID] - ρPerfectLocal1)] /.
    {Conjugate[Sin[θ]] → Sin[θ], Conjugate[Cos[θ]] → Cos[θ]}] /.
    {Conjugate[Sin[θ]] → Sin[θ], Conjugate[Cos[θ]] → Cos[θ]}]]
 $\left( \frac{\delta}{4} - 2 \delta \epsilon 3 + 8 \epsilon 3^2 + 6 \delta \epsilon 3^2 - 48 \epsilon 3^3 - 8 \delta \epsilon 3^3 + 136 \epsilon 3^4 + 4 \delta \epsilon 3^4 - \right.$ 
  224 ε35 + 224 ε36 - 128 ε37 + 32 ε38 + εk2 - δ εk2 - 8 ε3 εk2 + 8 δ ε3 εk2 +
  24 ε32 εk2 - 24 δ ε32 εk2 - 32 ε33 εk2 + 32 δ ε33 εk2 + 16 ε34 εk2 - 16 δ ε34 εk2  $\left. \right) /$ 
 $\left( 1 - 448 \epsilon 3^5 + 448 \epsilon 3^6 - 256 \epsilon 3^7 + 64 \epsilon 3^8 - \frac{1}{2} \delta (1 - 2 \epsilon 3)^4 (1 - 2 \epsilon k)^2 - \right.$ 
  2 εk + 2 εk2 - 64 ε33 (2 - εk + εk2) + 32 ε34 (9 - εk + εk2) -
  8 ε3 (1 - 2 εk + 2 εk2) + 8 ε32 (5 - 6 εk + 6 εk2)  $\left. \right)$ 

```

Error probability on a single output qubit.

```

Expand[Psuc[ε, ε, 0]]
Series[εout[ε, ε, 0], {ε, 0, 10}]
1 - 10 ε + 58 ε2 - 192 ε3 + 400 ε4 - 544 ε5 + 480 ε6 - 256 ε7 + 64 ε8
9 ε2 + 34 ε3 - 22 ε4 - 720 ε5 - 2756 ε6 + 1144 ε7 + 56 056 ε8 + 227 072 ε9 - 31 472 ε10 + O[ε]11

```

Above equations match the results of MEK

```

Normal[Series[Psuc[ε3 * t, ek * t, δ * t], {t, 0, 1}]] /. {t → 1}
Normal[Series[εout[ε3 * t, ek * t, δ * t2], {t, 0, 2}]] /. {t → 1}
1 -  $\frac{\delta}{2}$  - 8 ε3 - 2 εk
 $\frac{\delta}{4} + 8 \epsilon 3^2 + \epsilon k^2$ 

```

The leading order P_{suc} and ϵ_{out} .